# High Performance Computing

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#### **Review Paper**

#### Two-Level Checkpoint/Restart Modeling for GPGPU

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AICCSA 2011

### GPGPU

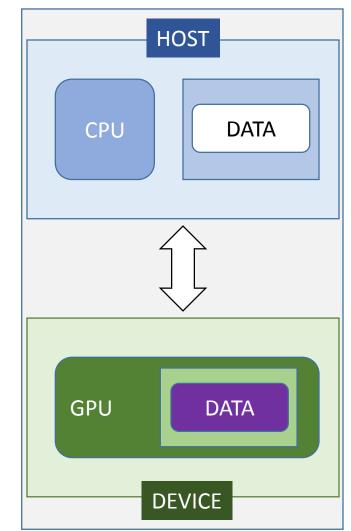
Nowadays, supercomputing applications has increasingly explored power of GPU and the cluster of GPUs for non-graphic applications [p276] →GPGPU

- GPU works as a co-processor of CPU
- Flow of execution of GPU application
- 1. Memory copy from Host to Device
- 2. Kernel execution
- 3. Memory copy from Device to Host
- CUDA
  - De facto standard for GPGPU
  - /usr/local/cuda on TSUBAME



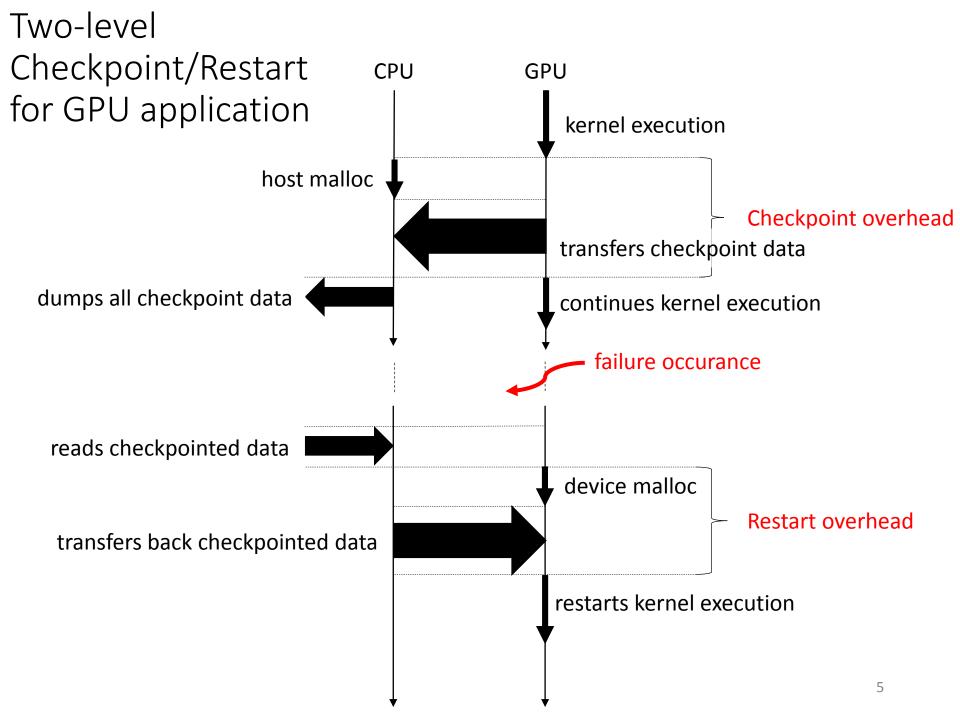






## Checkpoint/Restart for GPGPU

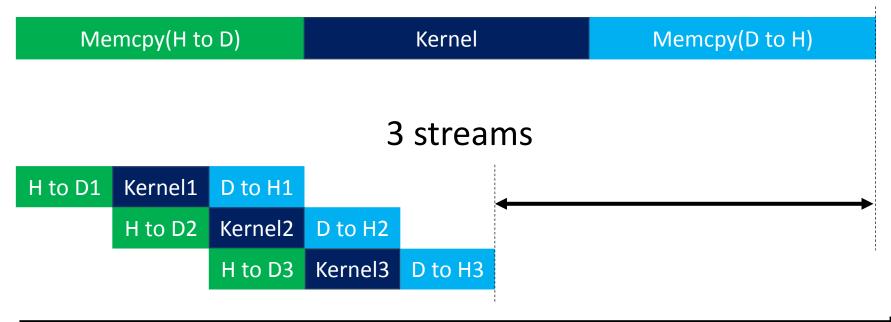
- BLCR[Paul, et al 2006]
  - Checkpoint/restart mechanism for Linux system
  - not works for GPU application
- VCCP[Hong, et al 2009]
  - Checkpoint/restart mechanism for virtual machines
  - may work
- NVCR[Nukada, et al 2011]
  - Transparent checkpoint/restart library for CUDA applications
  - not works for CUDA 4.0  $\sim$
- CheCL[Takizawa, et al 2011]
  - Transparent checkpoint/restart library for OpenCL applications
  - API proxy may be able to apply to CUDA



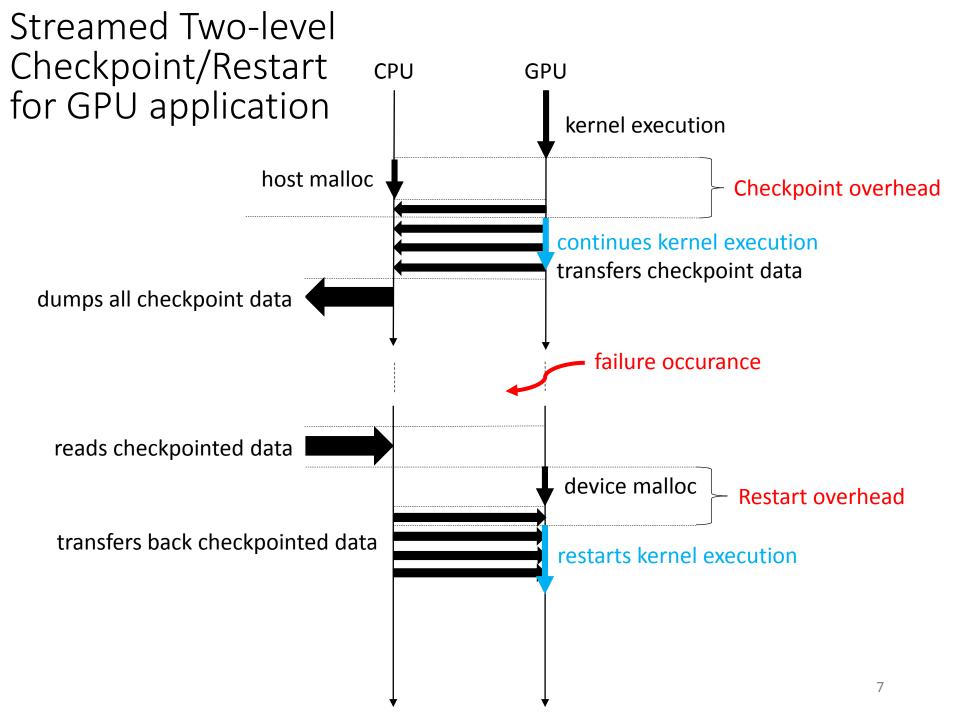
### CUDA Stream

• For overlap communication and computation

#### 1 stream(default)



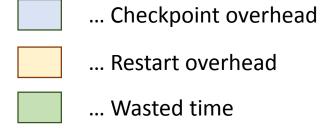
/usr/local/cuda/samples/0\_Simple/simpleStreams



### Experiment

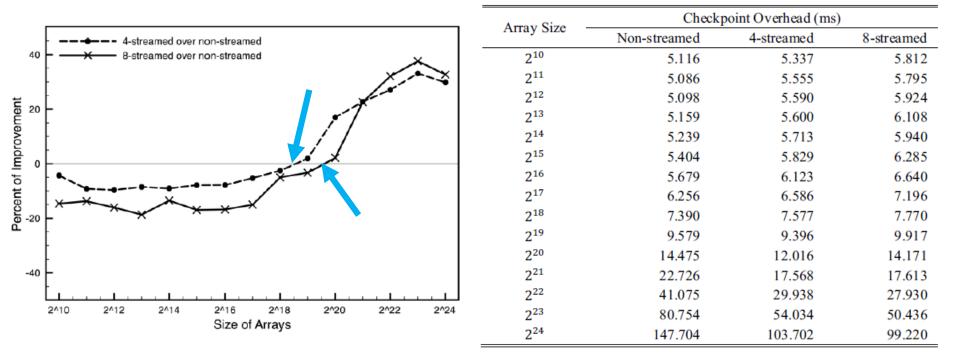
- NVIDIA GeForce GTX 295
  - compute capatibility 1.3
  - maximum number of blocks in a grid is 65535
  - maximum number of threads in block is 512
- Array addition C = A + B
  Array size 2 ~ 2
- Compare non-streamed, 4-streamed and 8-streamed

## Experiment-1

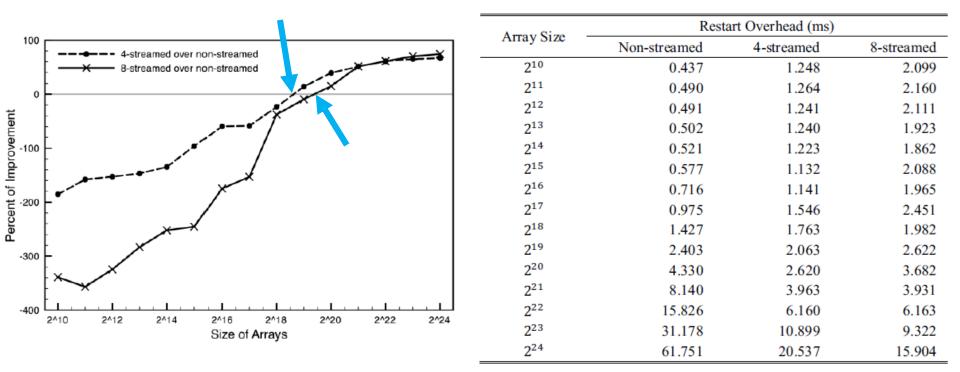


- Pseudo code
- 1. Execute kernel\_1() with L iterations
- 2. Synchronize all threads to prepare for memory copy
- 3. Allocate host memory for data checkpointing, i.e. array A, B, and C.
- 4. Do device-to-host memory copy of array A, B, and C.
- 5. Execute kernel\_2() with M iterations
- 6. Failure occurrence (Free all device memory)
- 7. Reallocate device memory for array A, B, and C.
- 8. Do host-to-device memory of array A, B, and C.
- 9. Re-execute kernel\_2() with M iterations.
- 10. Execute kernel\_3() with N iterations.
- 11. Do device-to-host memory copy of the result array C

#### Checkpoint overheads

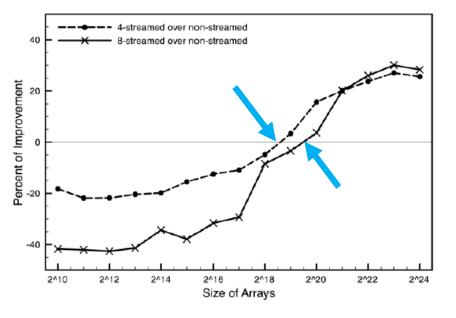


#### Restart overheads



#### Wasted times

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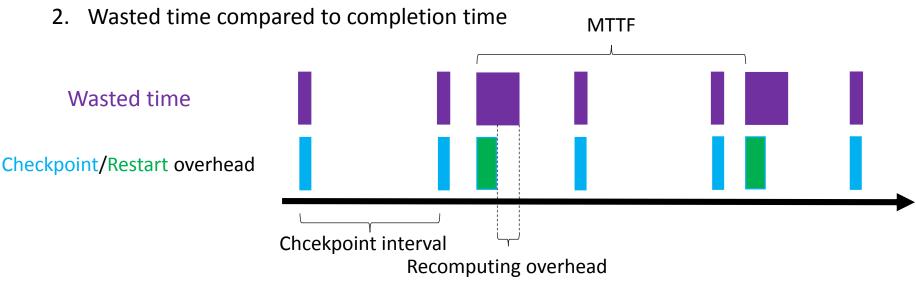
Array Size	Wasted Time (ms)		
	Non-streamed	4-streamed	8-streamed
210	5.665	6.698	8.025
211	5.688	6.931	8.067
212	5.701	6.943	8.147
2 <sup>13</sup>	5.776	6.954	8.146
2 <sup>14</sup>	5.936	7.112	7.978
2 <sup>15</sup>	6.322	7.302	8.714
216	6.964	7.834	9.175
2 <sup>17</sup>	8.255	9.156	10.671
2 <sup>18</sup>	10.803	11.326	11.738
2 <sup>19</sup>	15.913	15.389	16.469
2 <sup>20</sup>	26.573	22.404	25.620
2 <sup>21</sup>	46.352	37.016	37.029
222	87.749	66.946	64.941
2 <sup>23</sup>	173.624	126.625	121.450
224	332.888	247.673	238.557

## Experiment-2

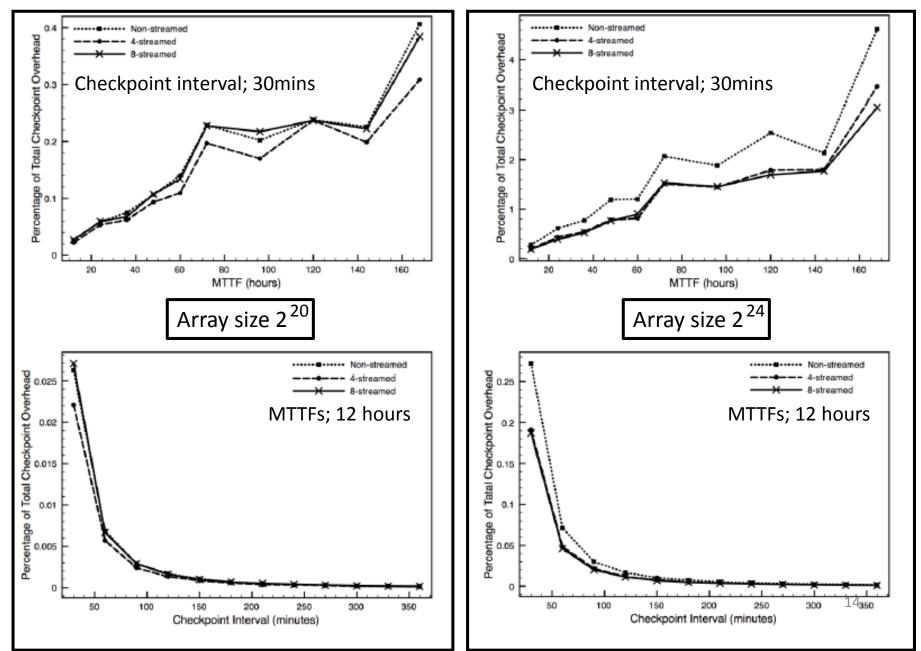
Consider mean-time-to-failures(MTTFs)

Size of array	2 <sup>20</sup> and 2 <sup>24</sup>	
MTTFs	12 hours to 7 days	
Checkpoint interval	30 and 120 mins	

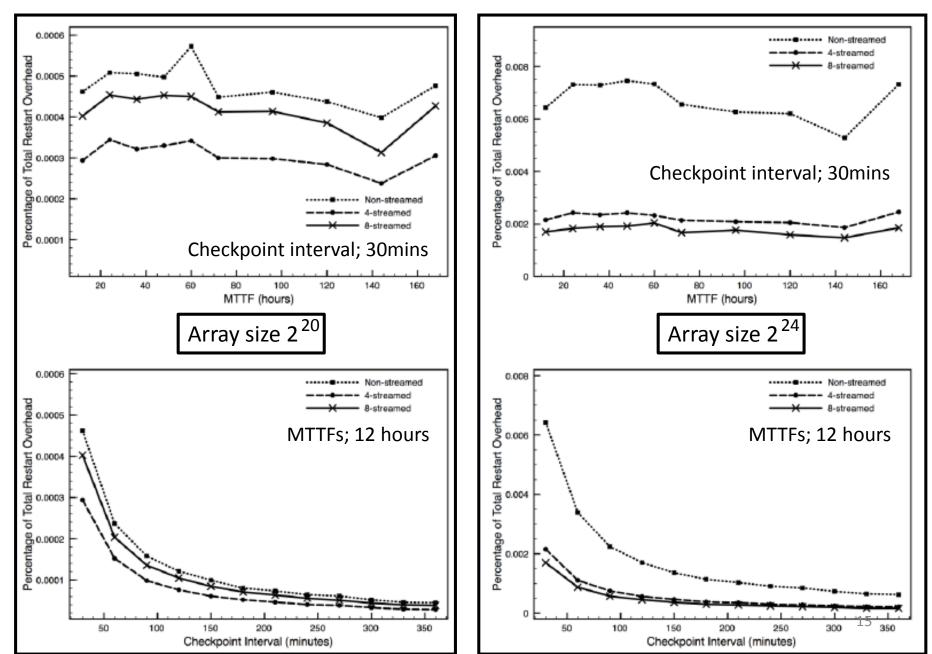
1. Checkpoint/restart overheads compared to wasted time



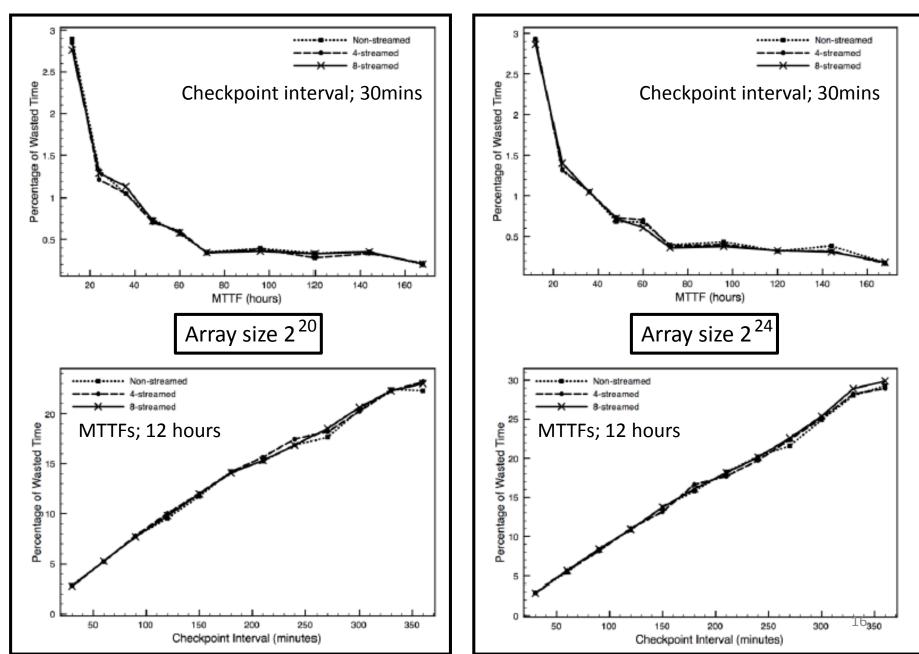
#### Checkpoint overhead/wasted time



#### Restart overhead/wasted time



### Wasted time/completion time



## Conclusion

- GPGPU
  - GPUs for non-graphic applications
  - De facto standard; CUDA
- Checkpoint/restart modeling for GPGPU
  - Two-level
  - Streamed
- Streamed Checkpoint/Restart model has advantage over non-streamed model when data size is enough big

### Discussion

- How to implement the streamed model?
  - especially for an application with overlapped kernel?
- The results of experiments are affected by a characteristic of application.